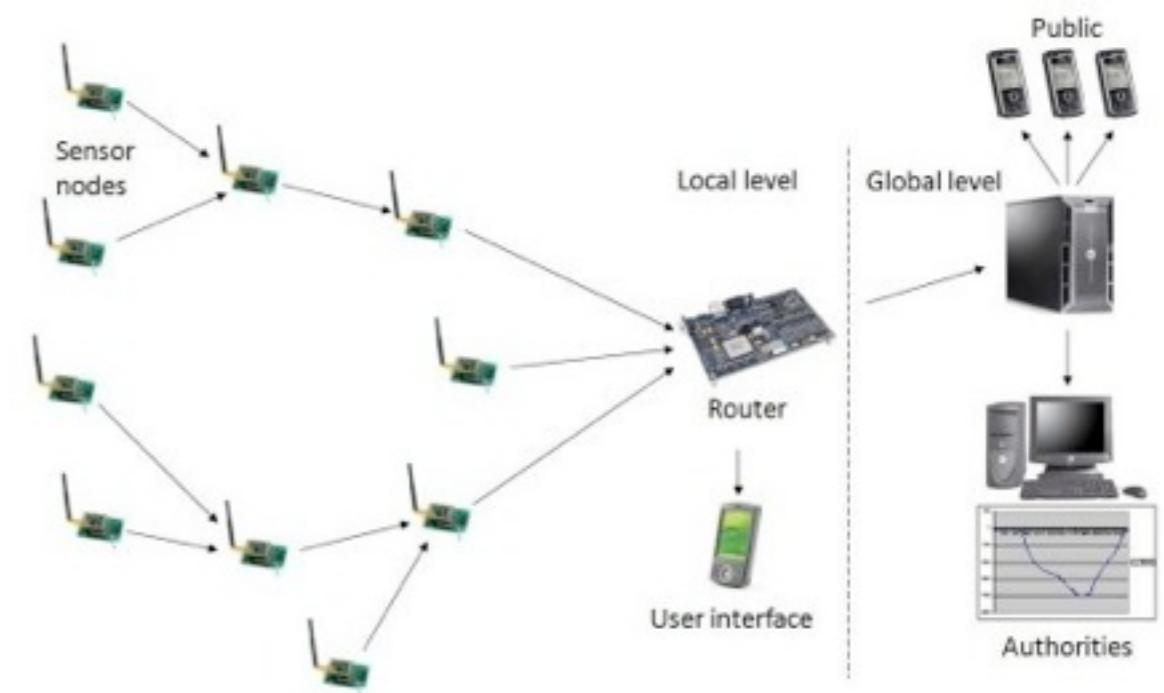
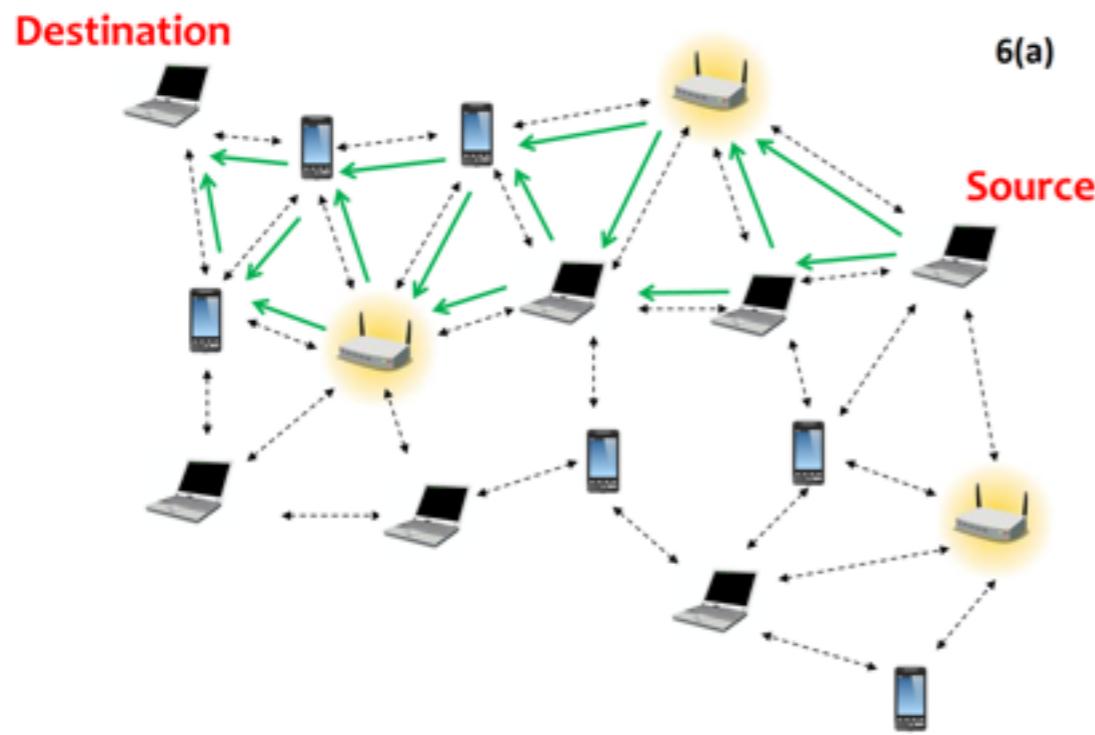


Trees & Routing (I)

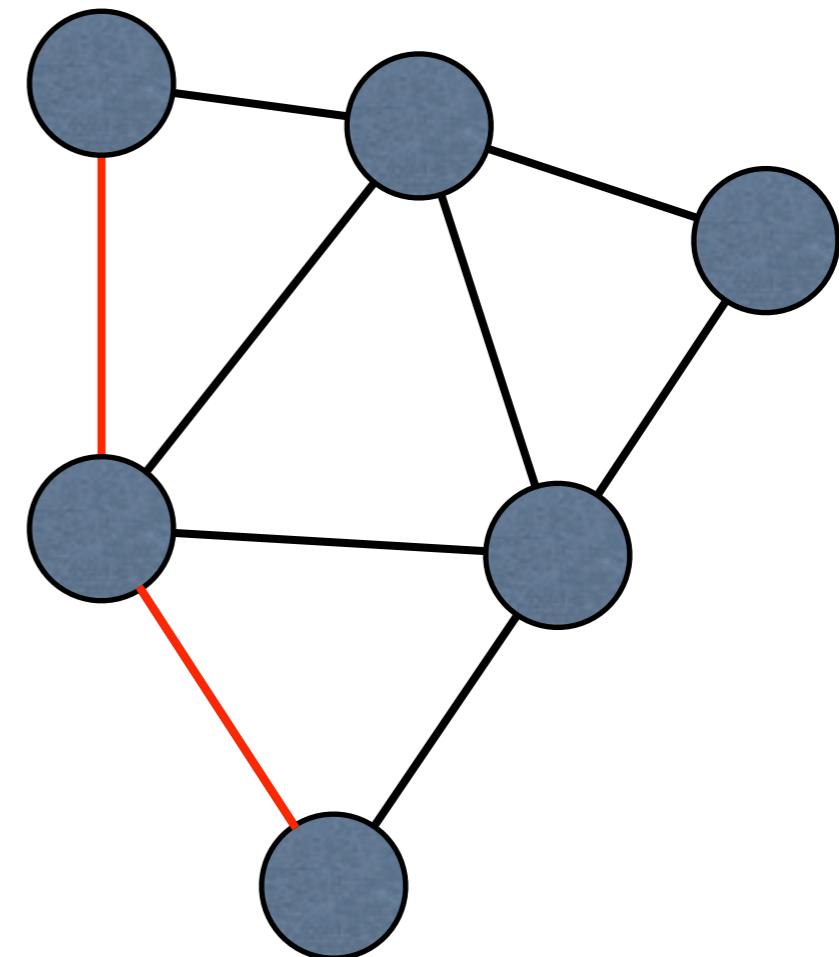
BFS

Why (spanning) trees ?



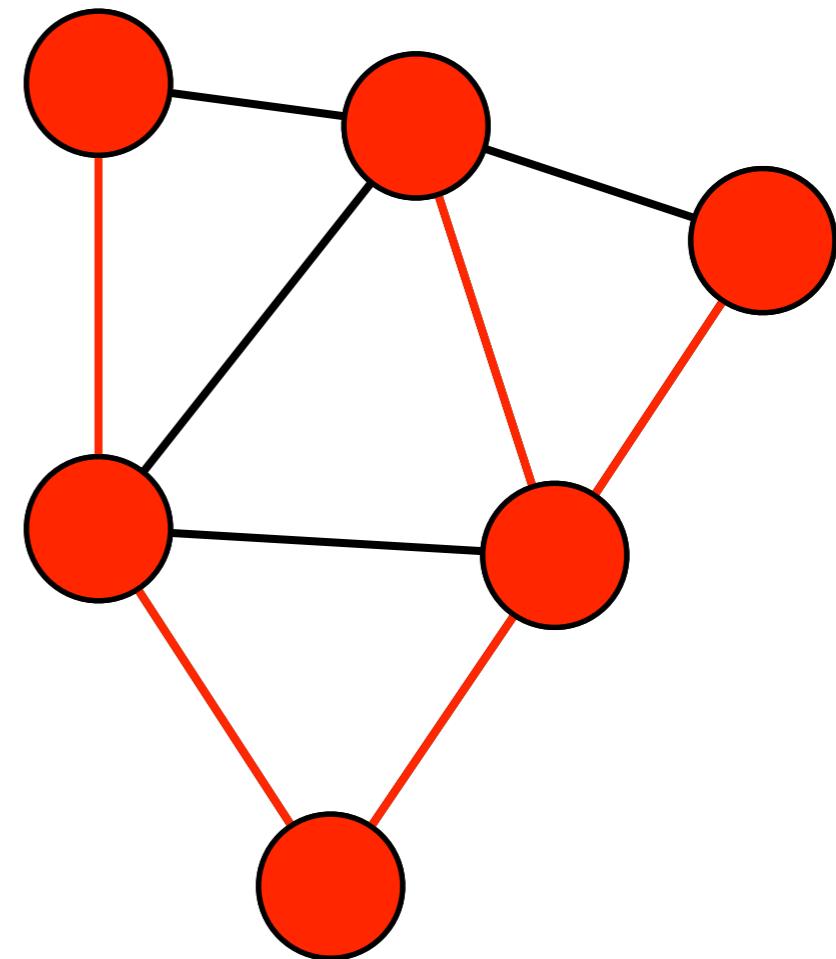
Spanning Tree

- Given a graph $G=(V,E)$, a **Spanning Tree** of G is an acyclic subgraph, spanning all the nodes in V .



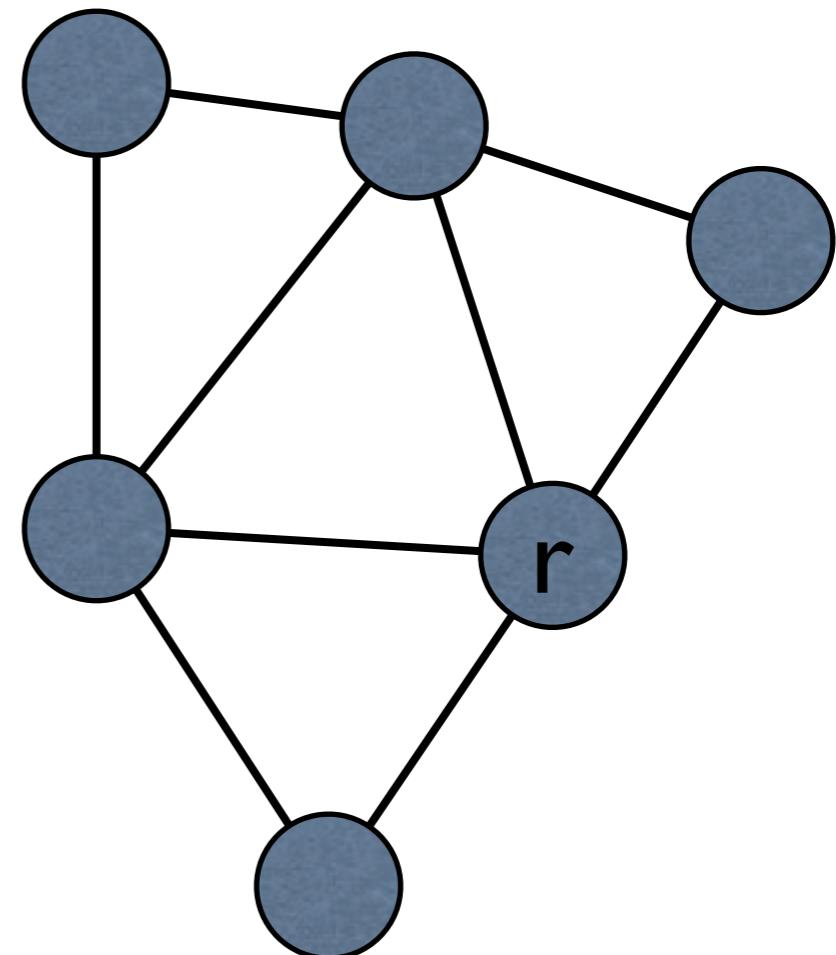
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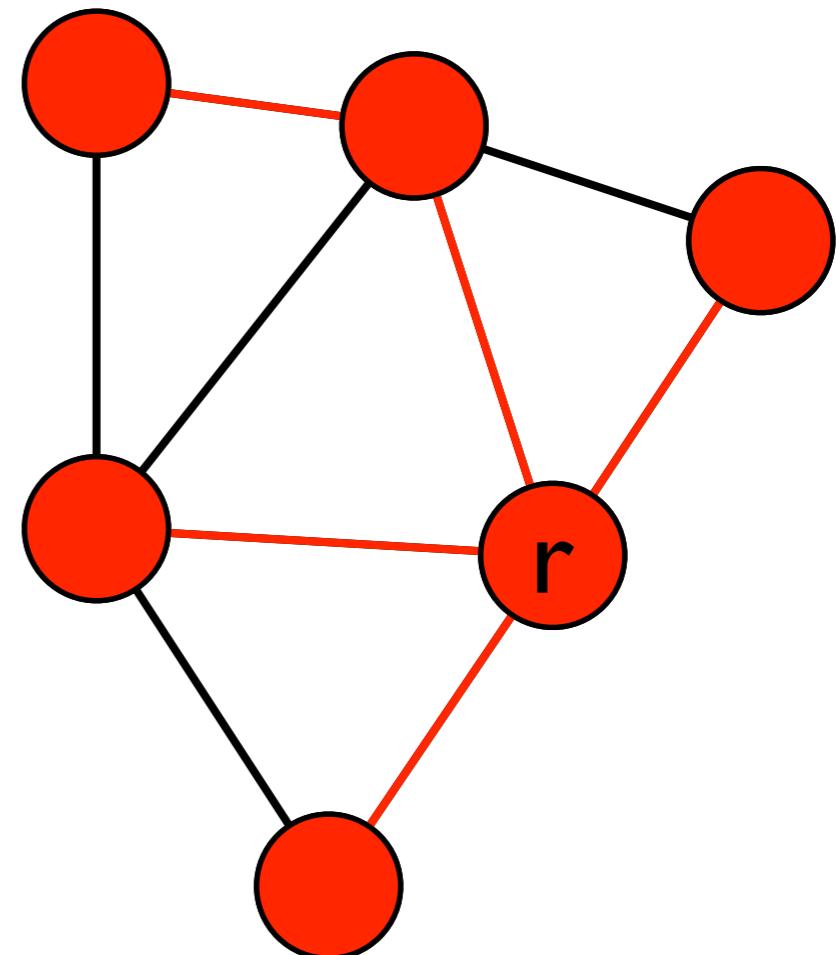
Breadth-first Spanning Tree (BFS)

- Given a graph $G=(V,E)$ and a node r in V , a **BFS** tree of G rooted at r is a spanning tree of G such that for any node v the path leading from the root to v in the tree is of the minimum length possible.



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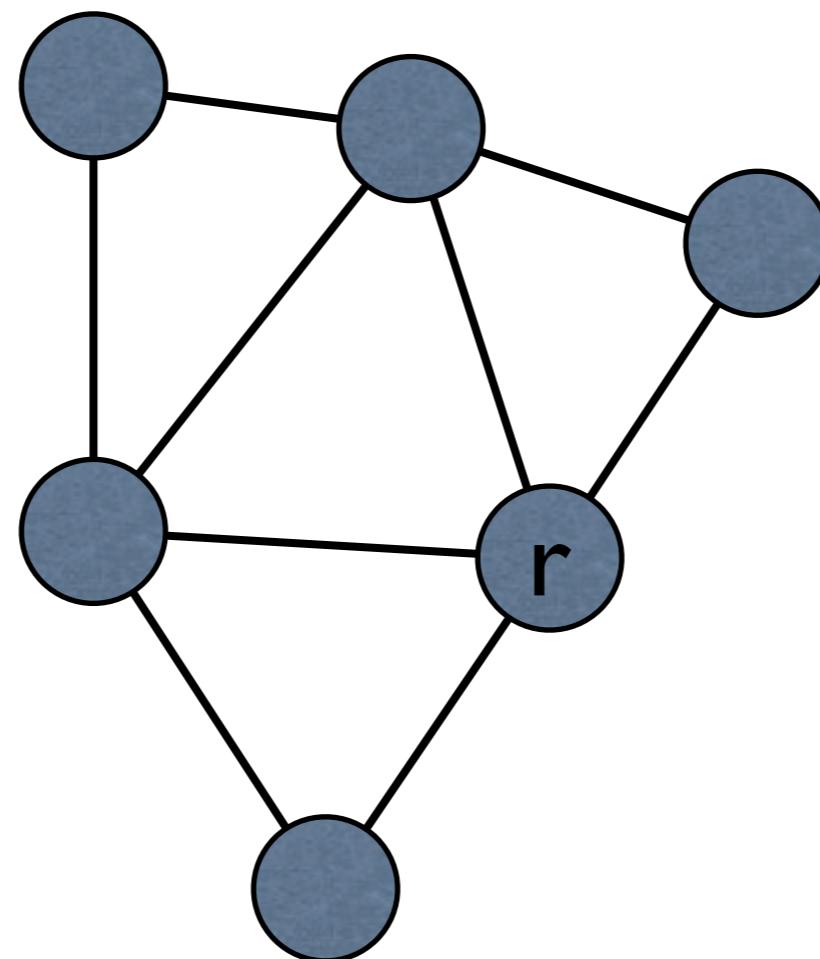


BFS Construction

BFS - strategies

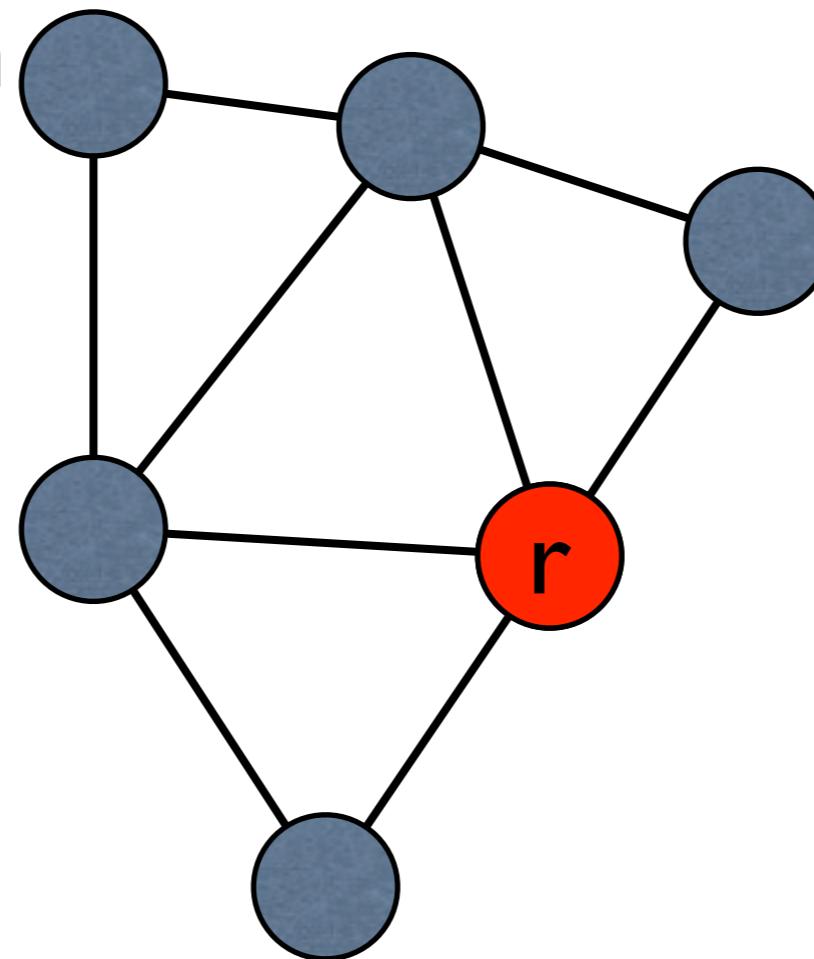
- «link state» (Dijkstra)
 - grow on the border
- «distance vector» (Bellman-Ford)
 - compute shortest paths

Dijkstra BFS

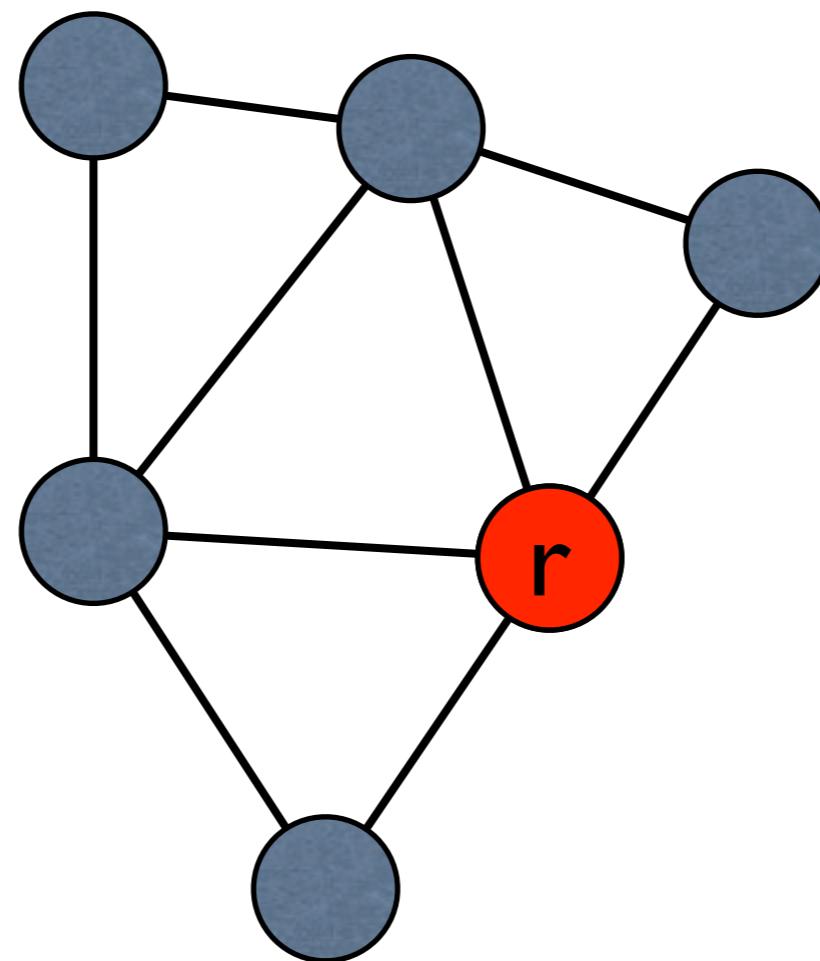


Dijkstra BFS

root, r , broadcasts in
the current tree, T ,
«pulse»

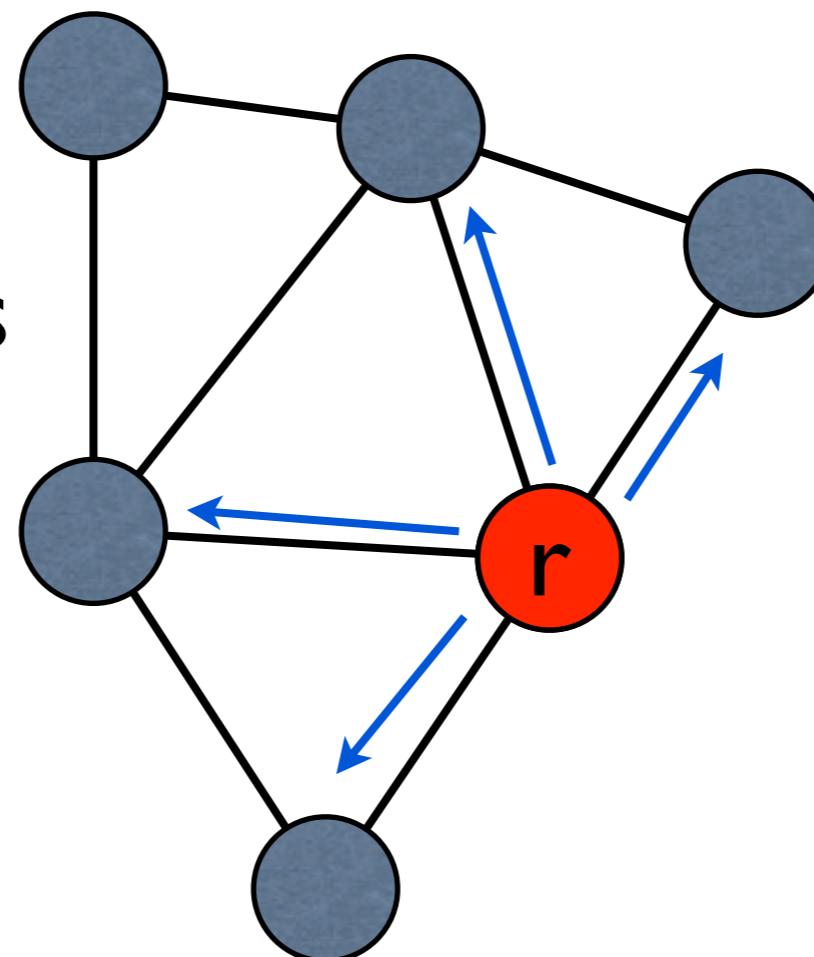


Dijkstra-like BFS

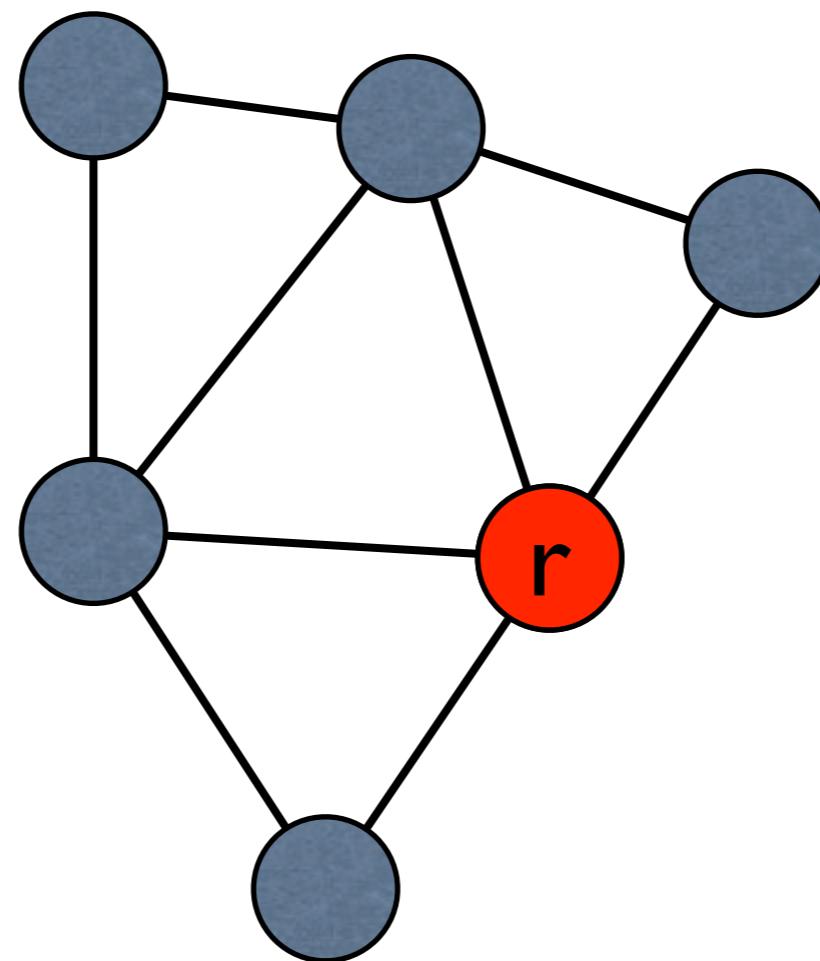


Dijkstra-like BFS

upon reception of
«pulse», leafs of the
current tree send
«layer» to neighbors
(except to parents)

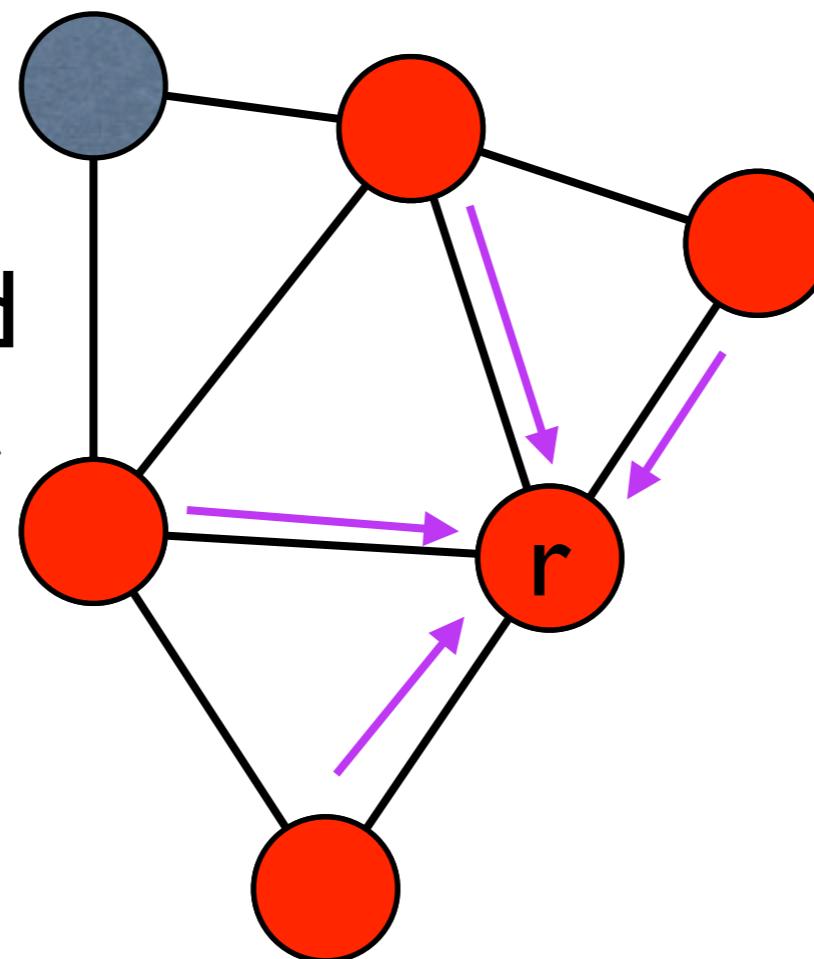


Dijkstra-like BFS



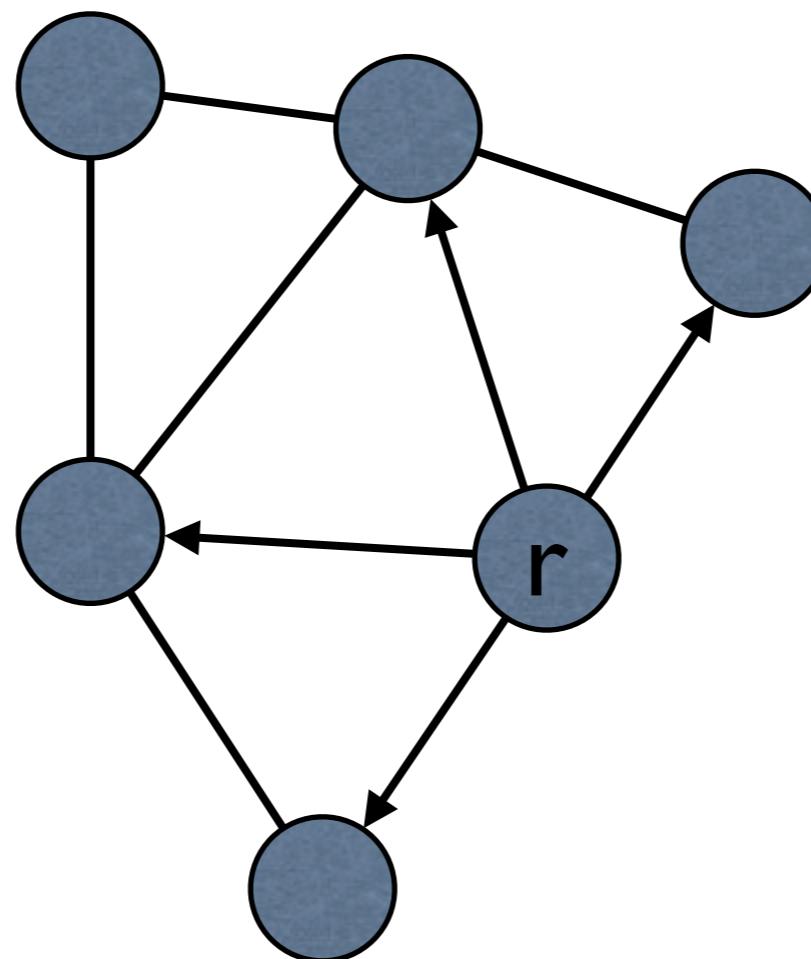
Dijkstra-like BFS

upon reception of
the first «layer» a
node chooses the
sender as parent and
sends back an «ack»



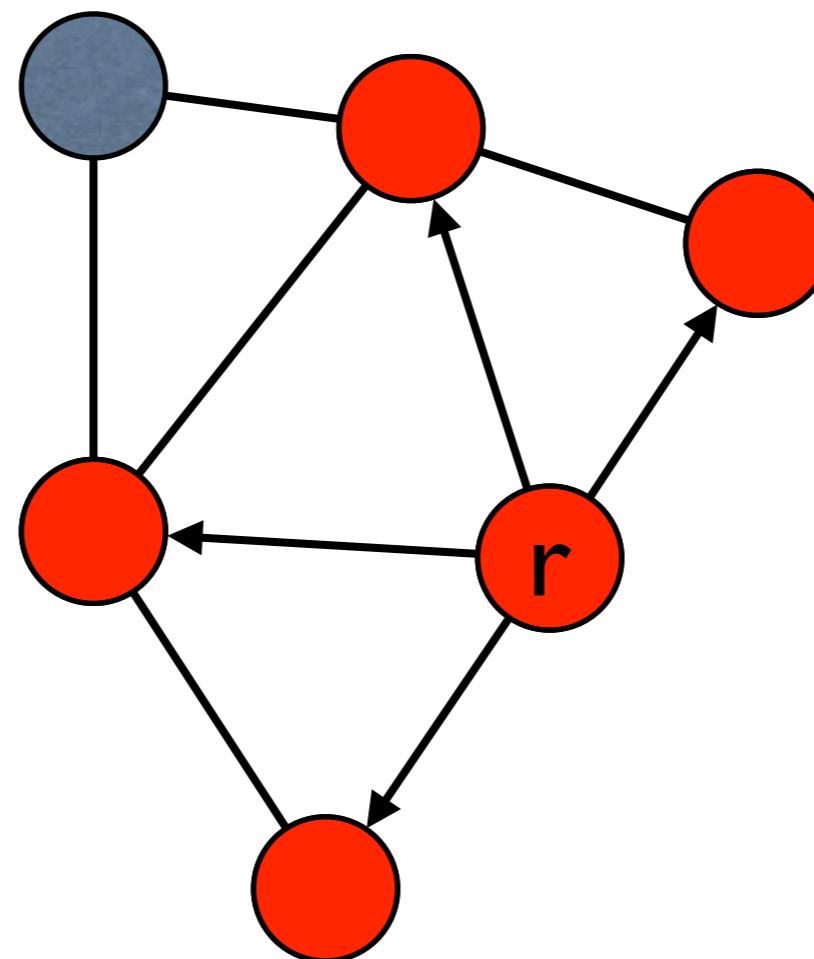
Dijkstra-like BFS

upon reception of
an «**ack**» a node
register the sender
as its child

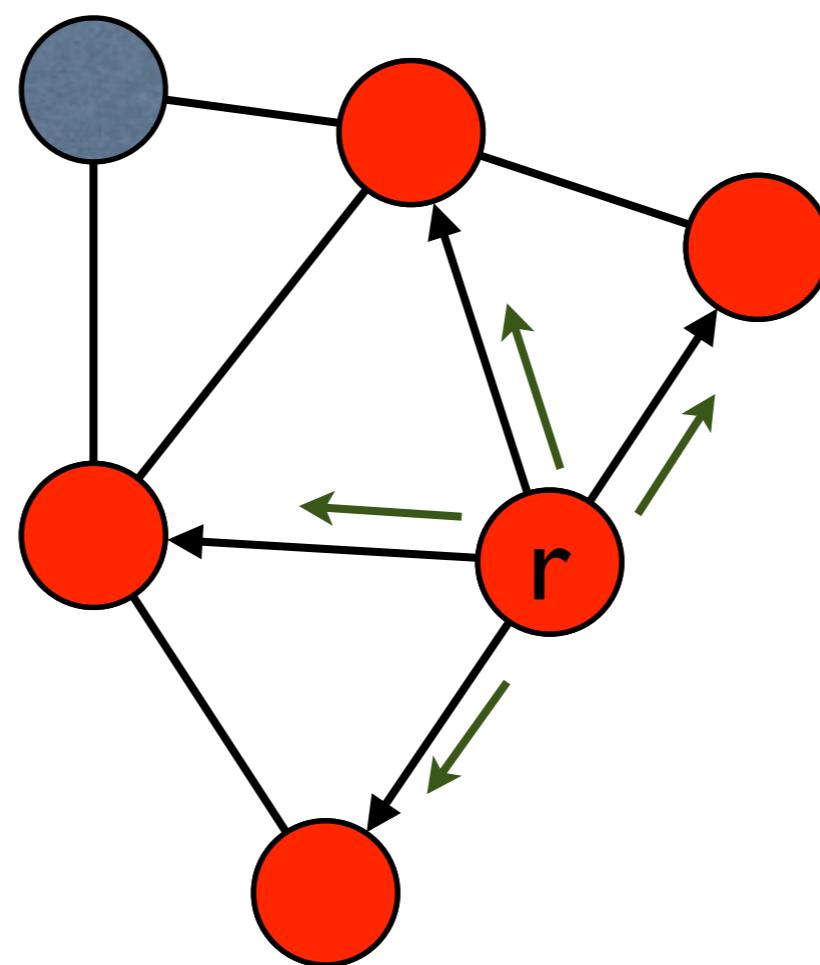


Dijkstra-like BFS

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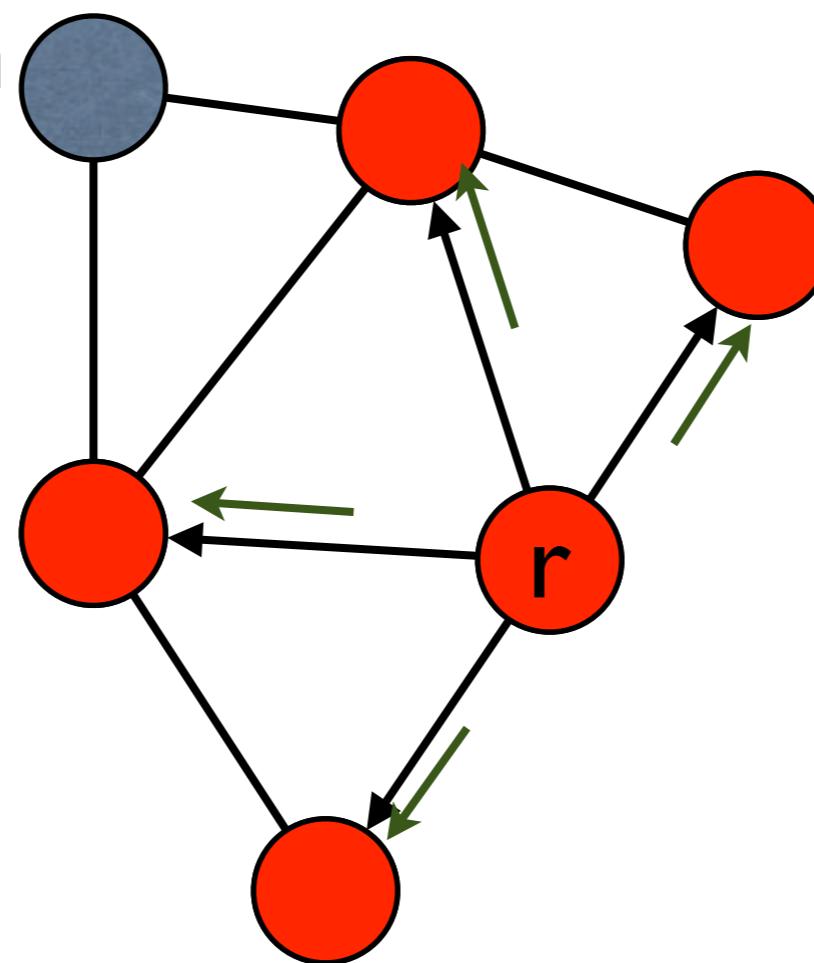


Dijkstra-like BFS

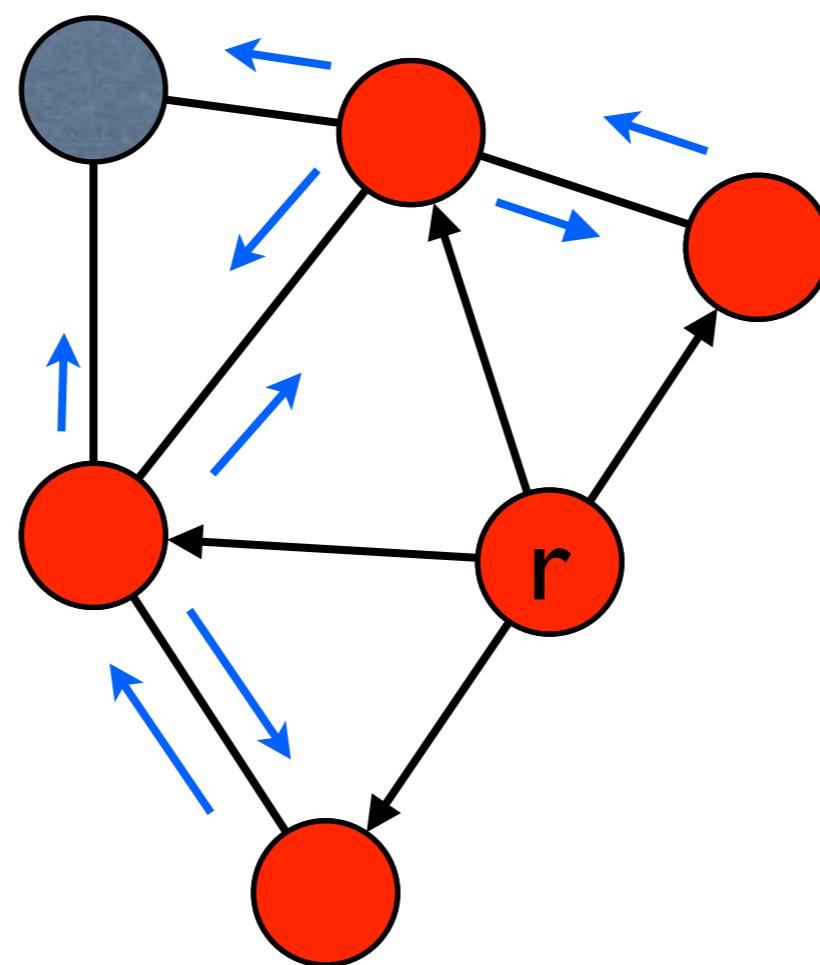


Dijkstra-like BFS

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the current tree, T ,
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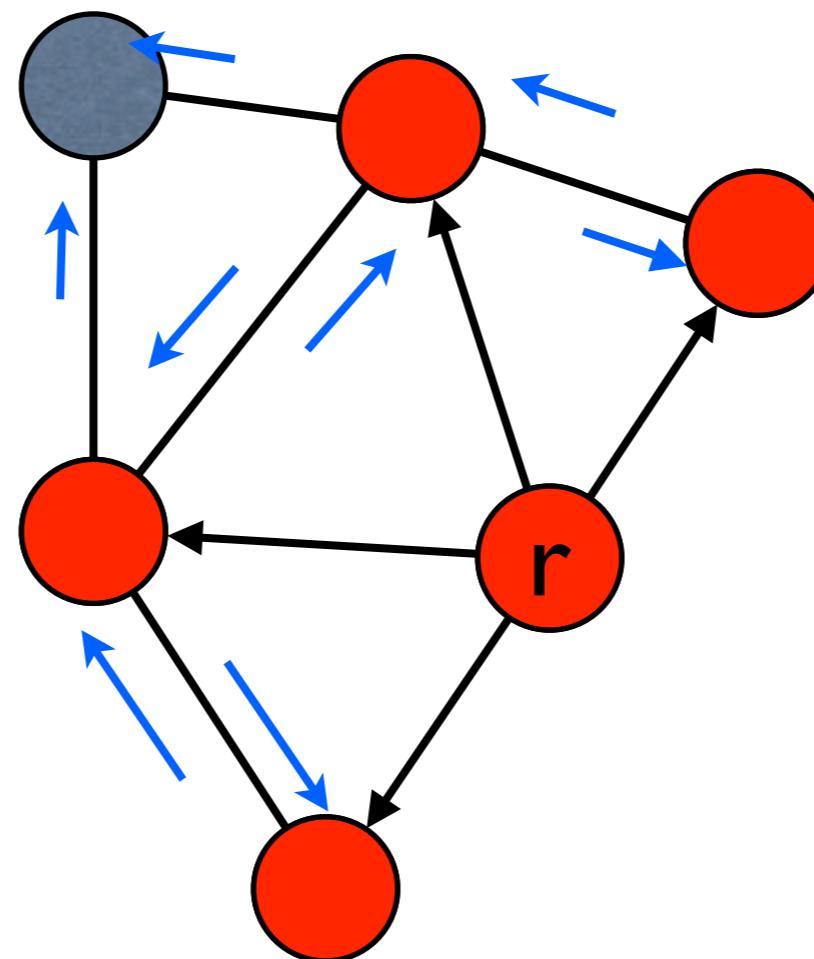


Dijkstra-like BFS



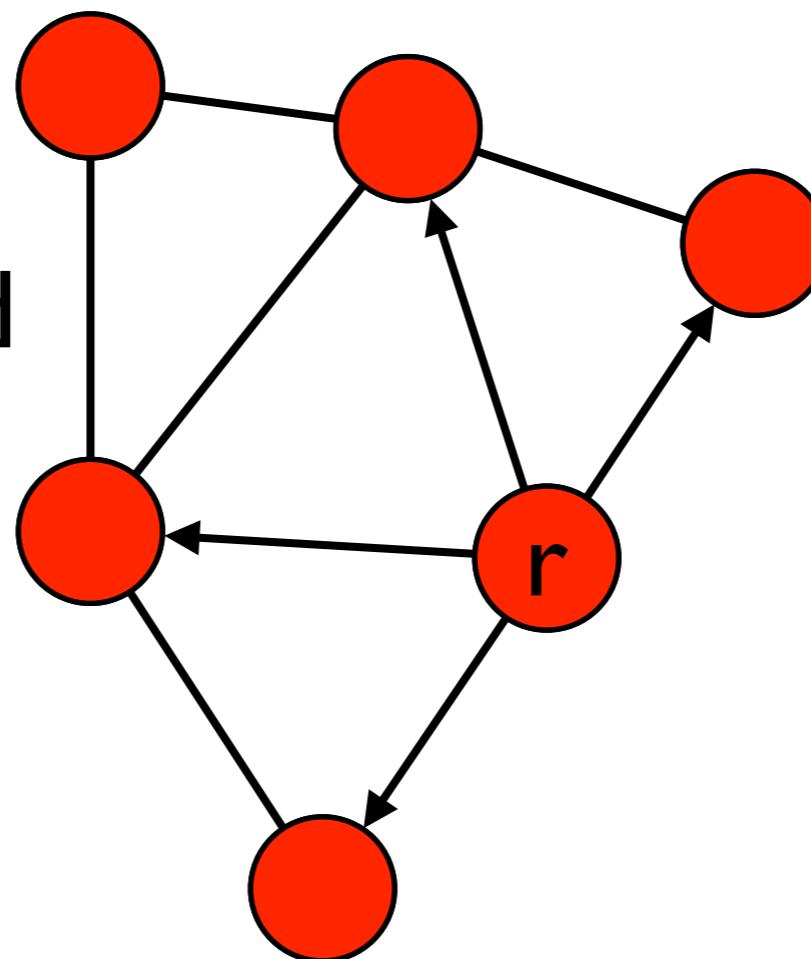
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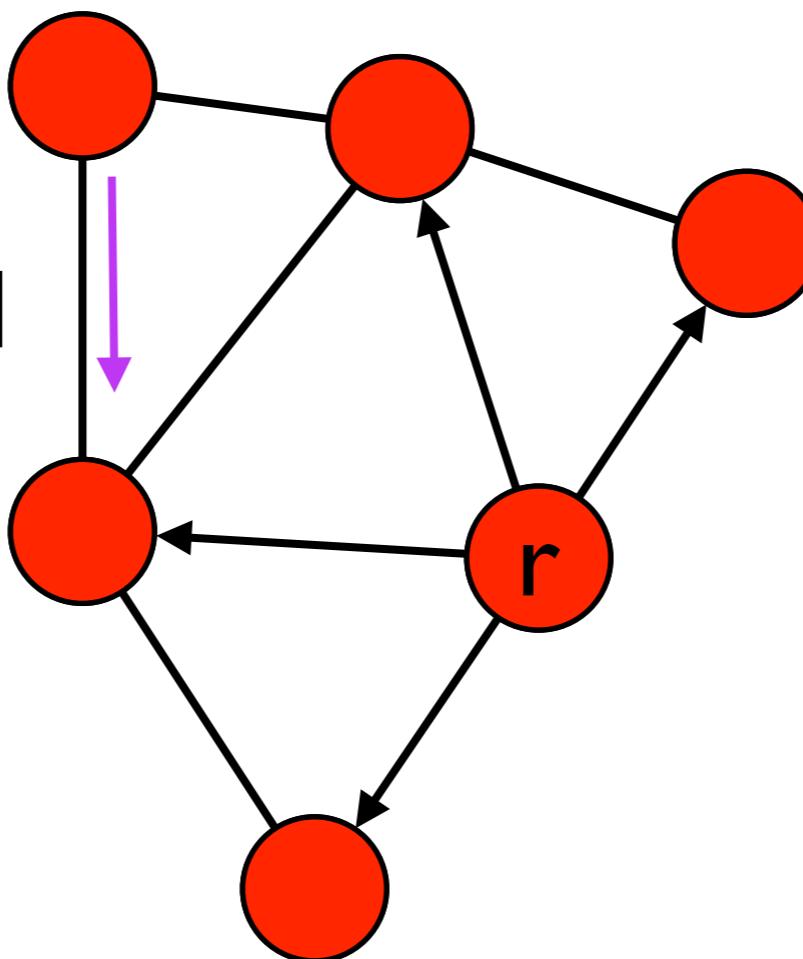
Dijkstra-like BFS

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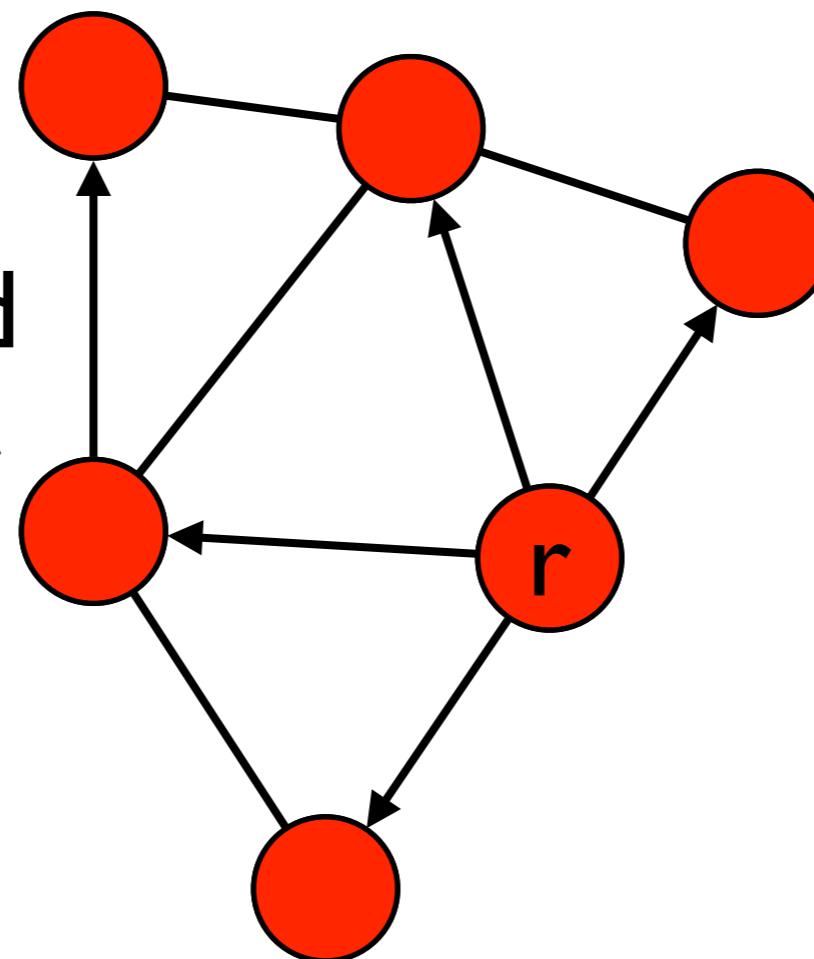
Dijkstra-like BFS

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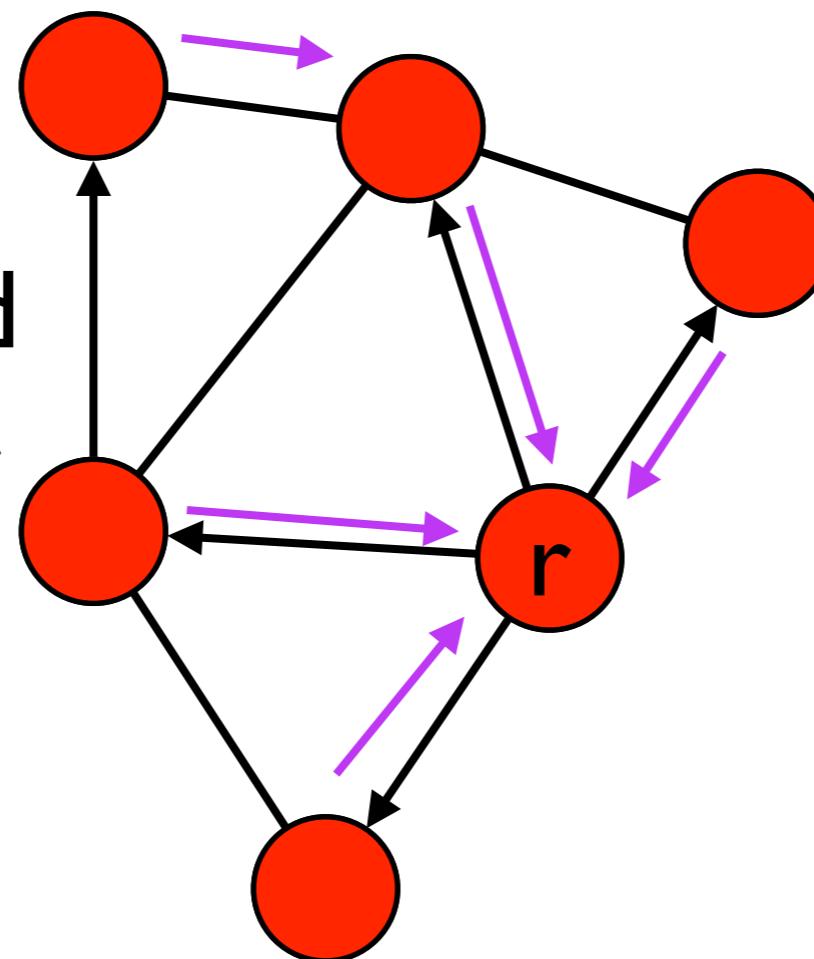
Dijkstra-like BFS

upon reception of
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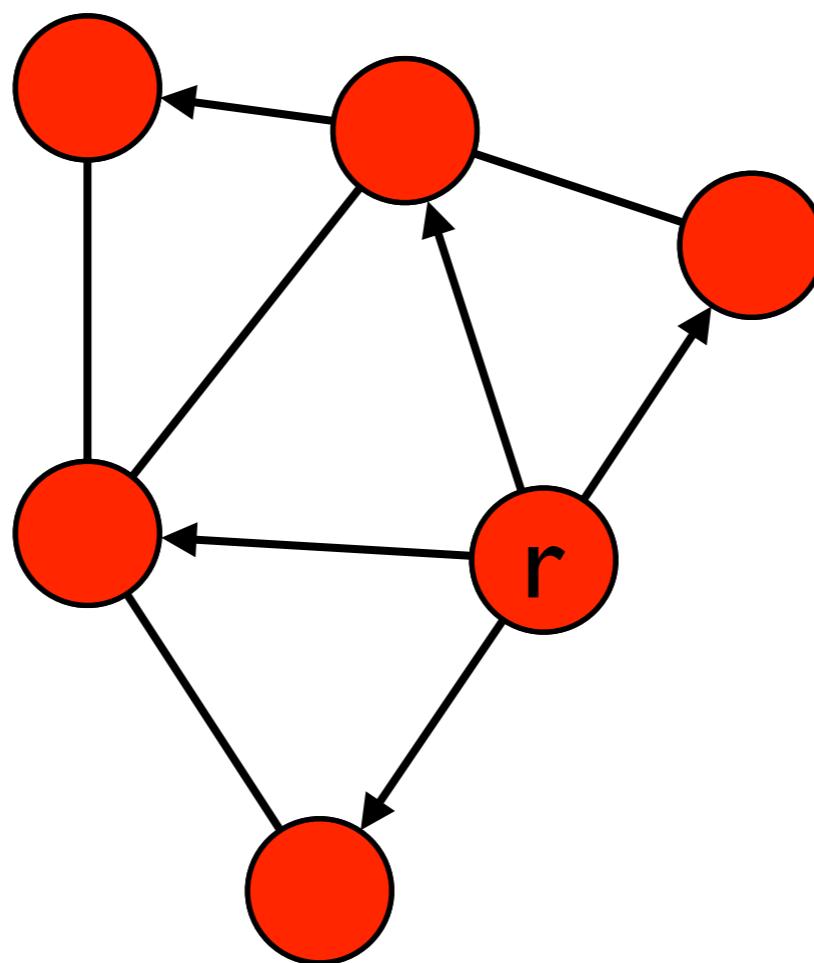


Dijkstra-like BFS

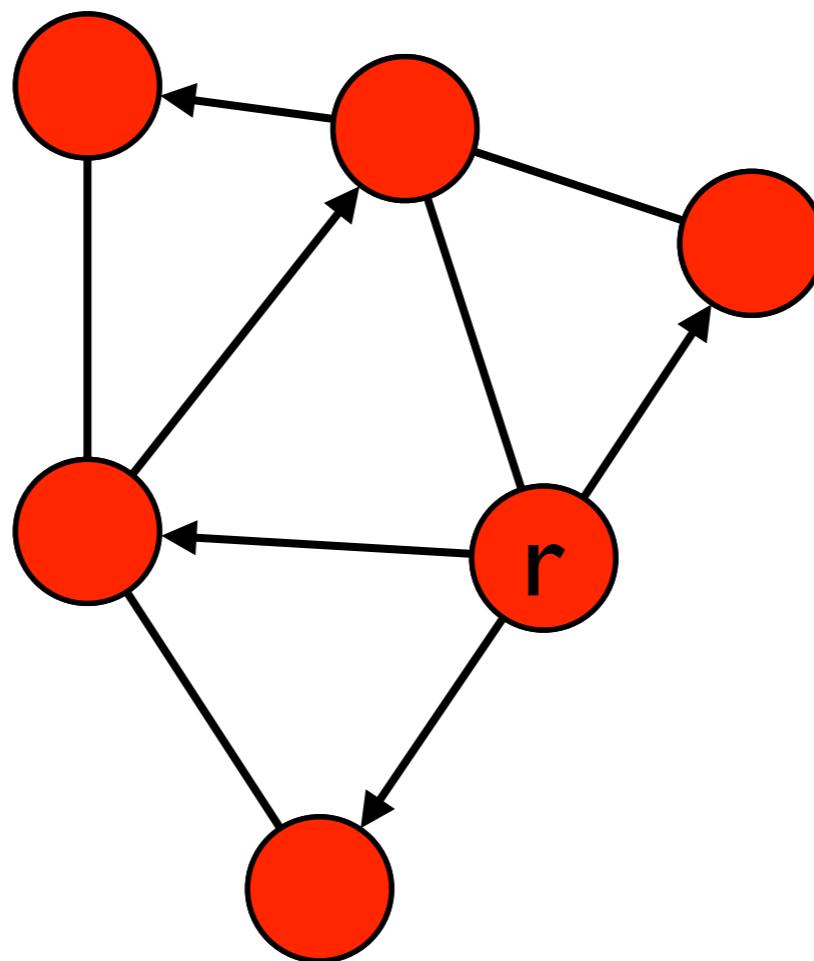
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Dijkstra-like BFS ?



Dijkstra-like BFS ?



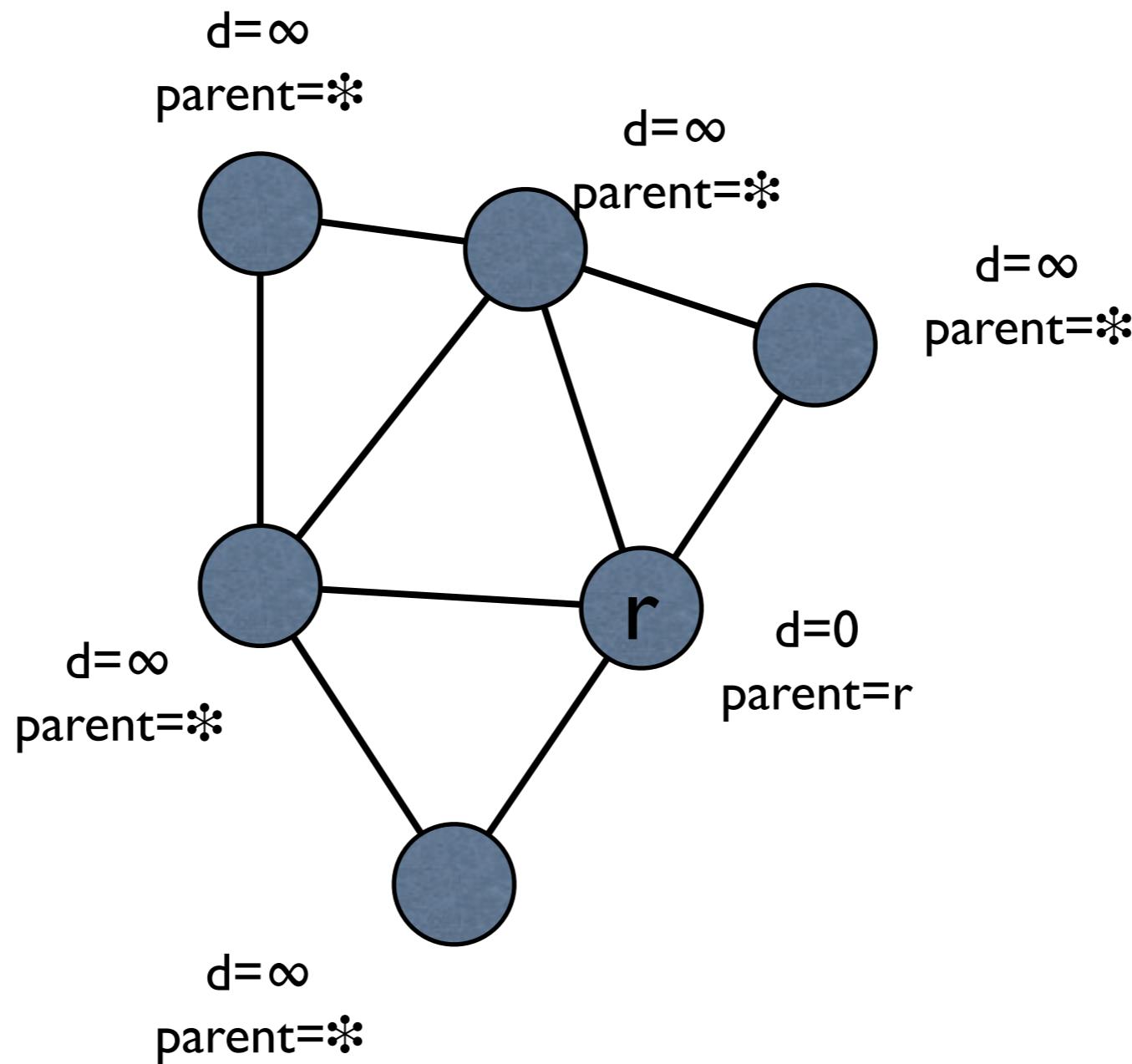
Dijkstra-like BFS - Complexity

- Time Complexity: $O(\text{Diam}^2(G))$
- Message Complexity: $O(n\text{Diam}(G) + |E|)$

Bellman-Ford-like BFS

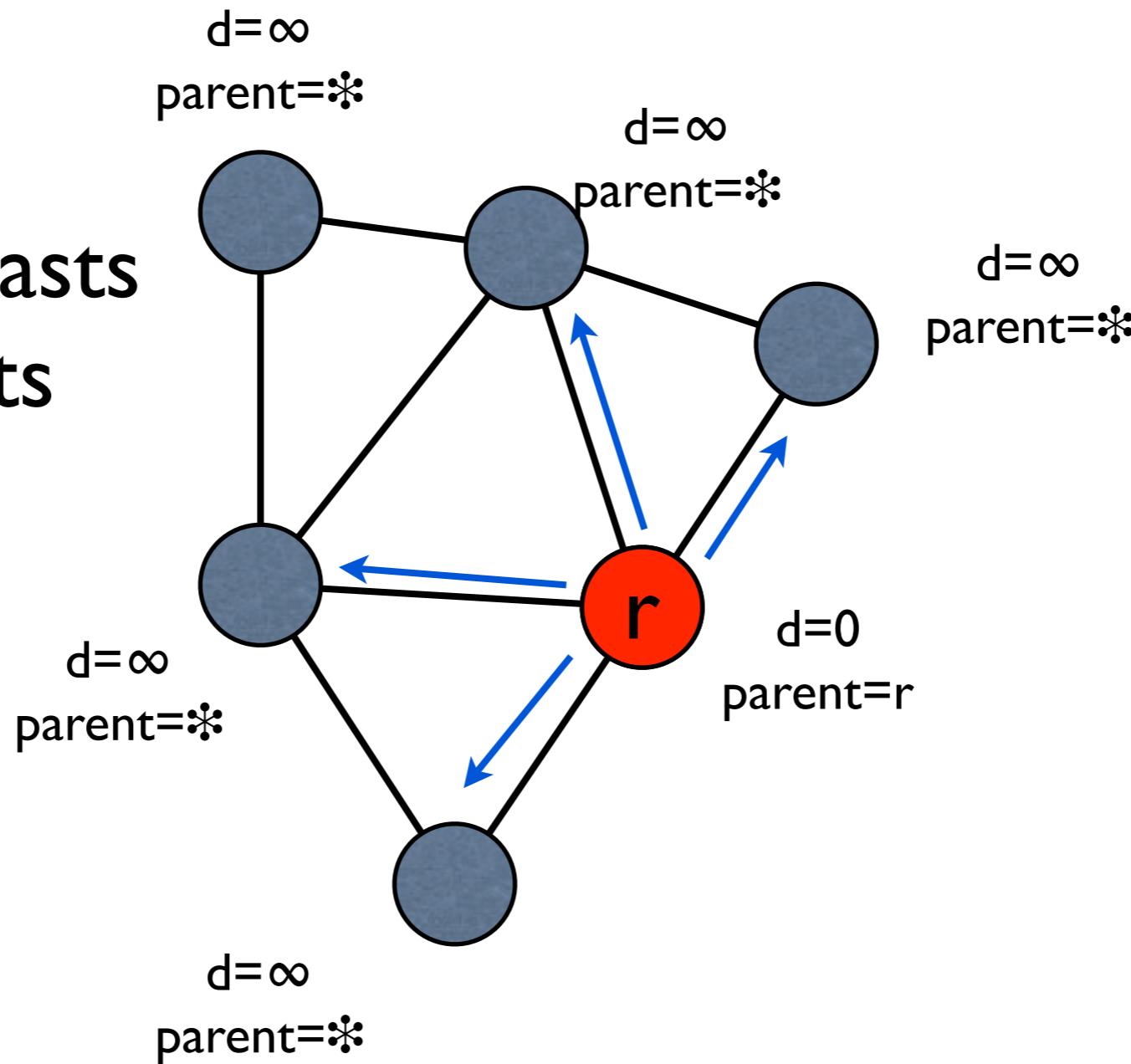
- each node maintains :
 - its shortest distance to the root d (initially $d=\infty$)
 - its parent to the root on the shortest distance

Bellman-Ford-like BFS

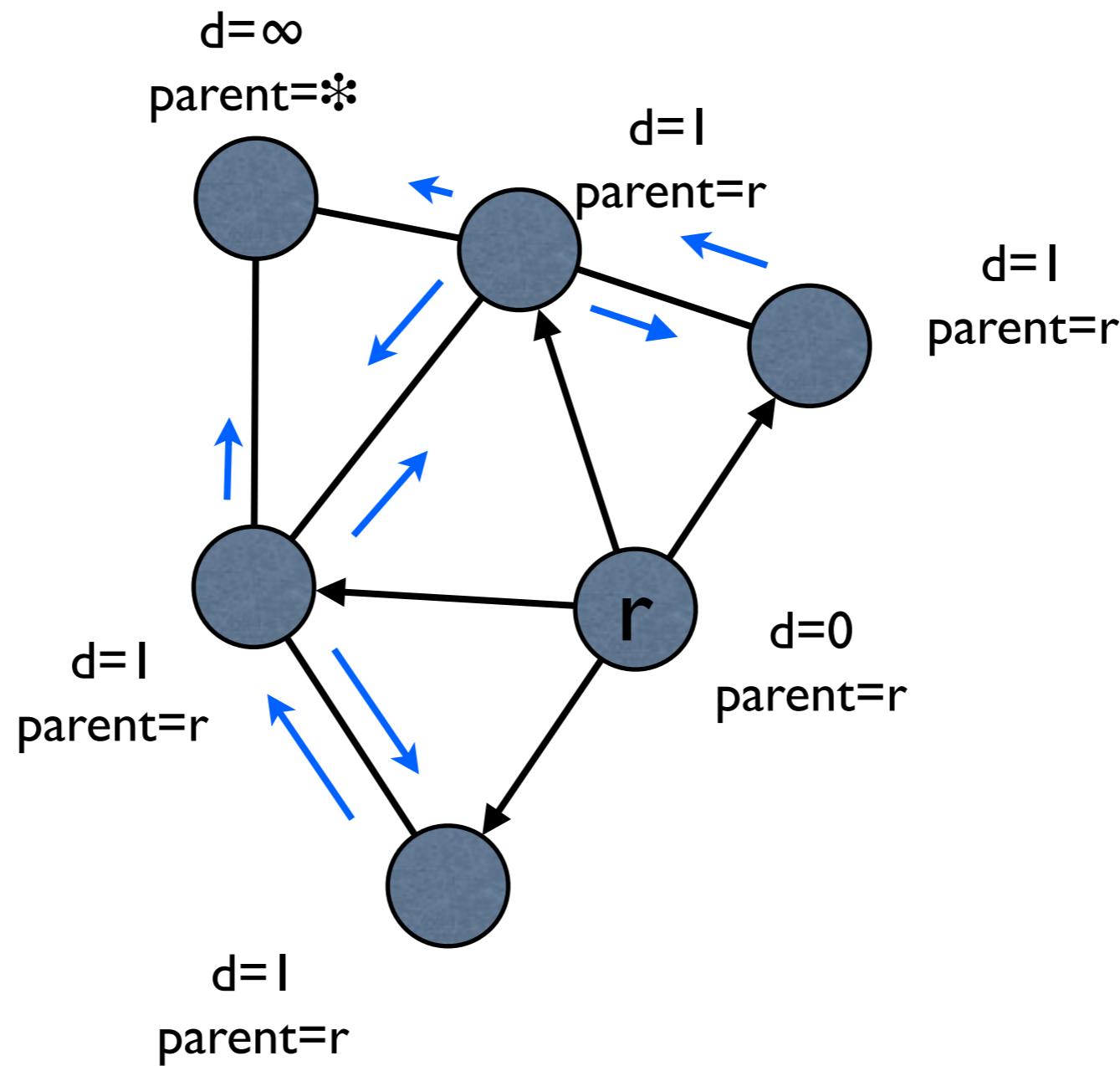


Bellman-Ford-like BFS

root, r , broadcasts
 $d=0$ to all its
neighbors

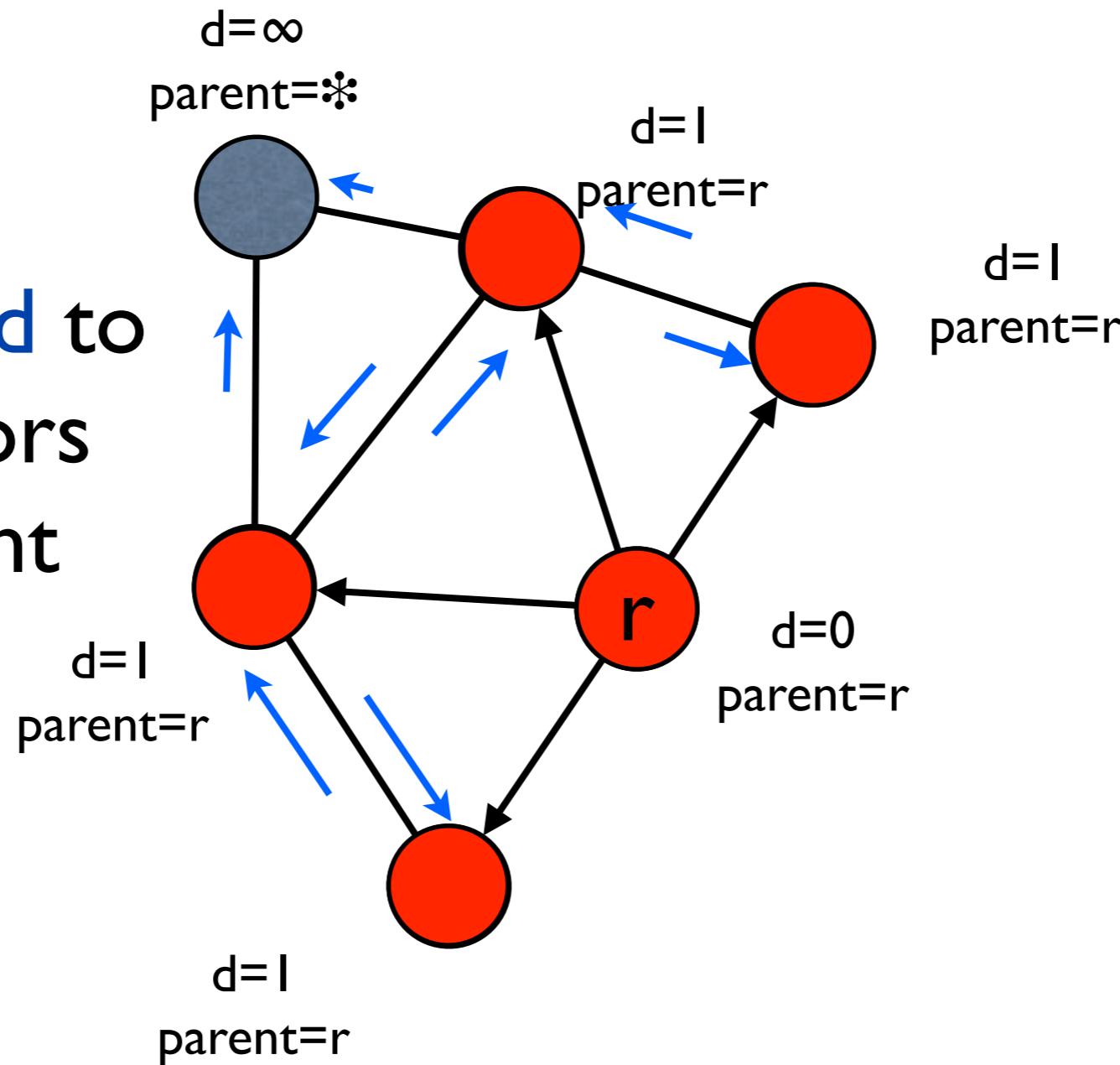


Bellman-Ford-like BFS

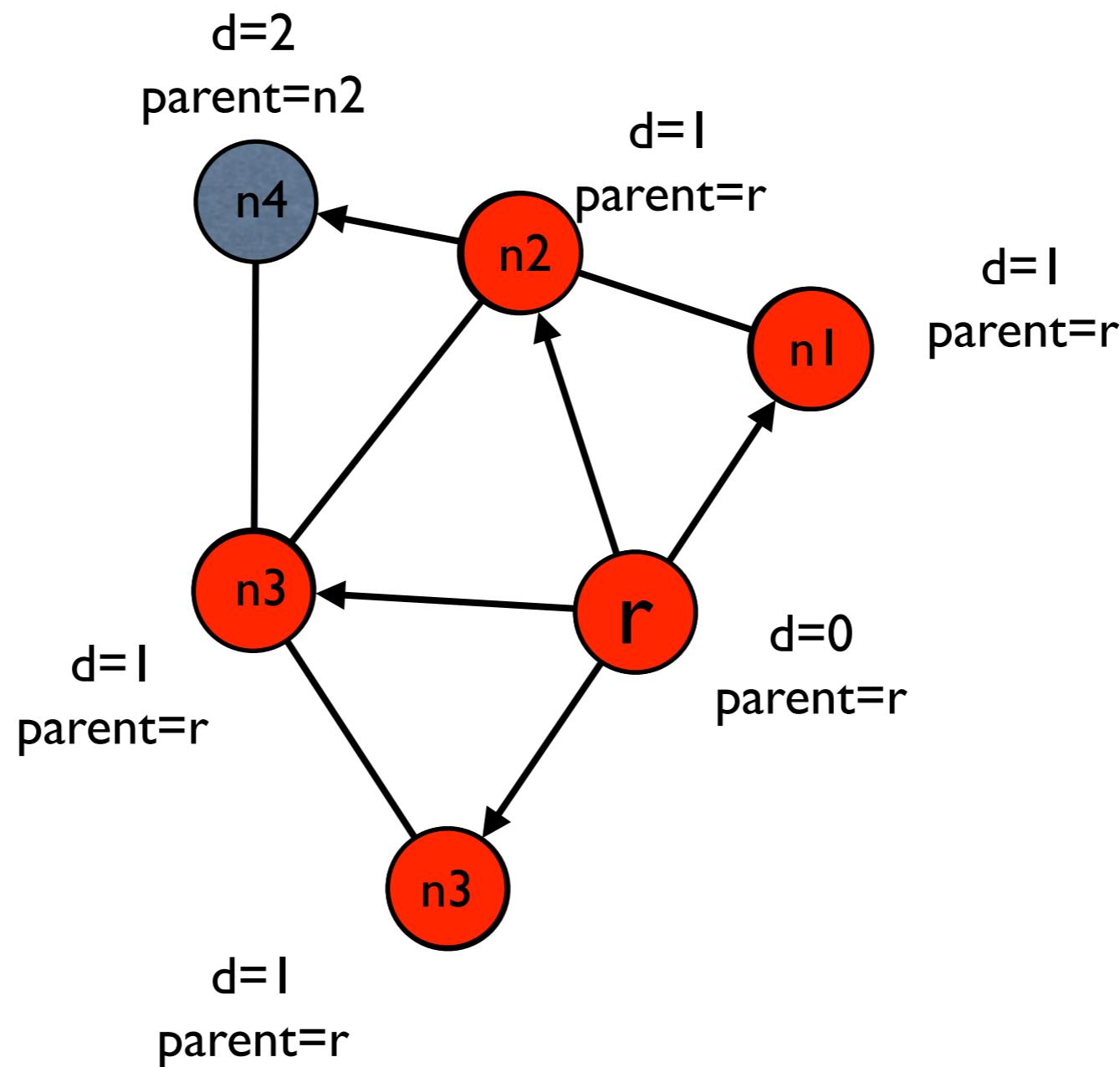


Bellman-Ford-like BFS

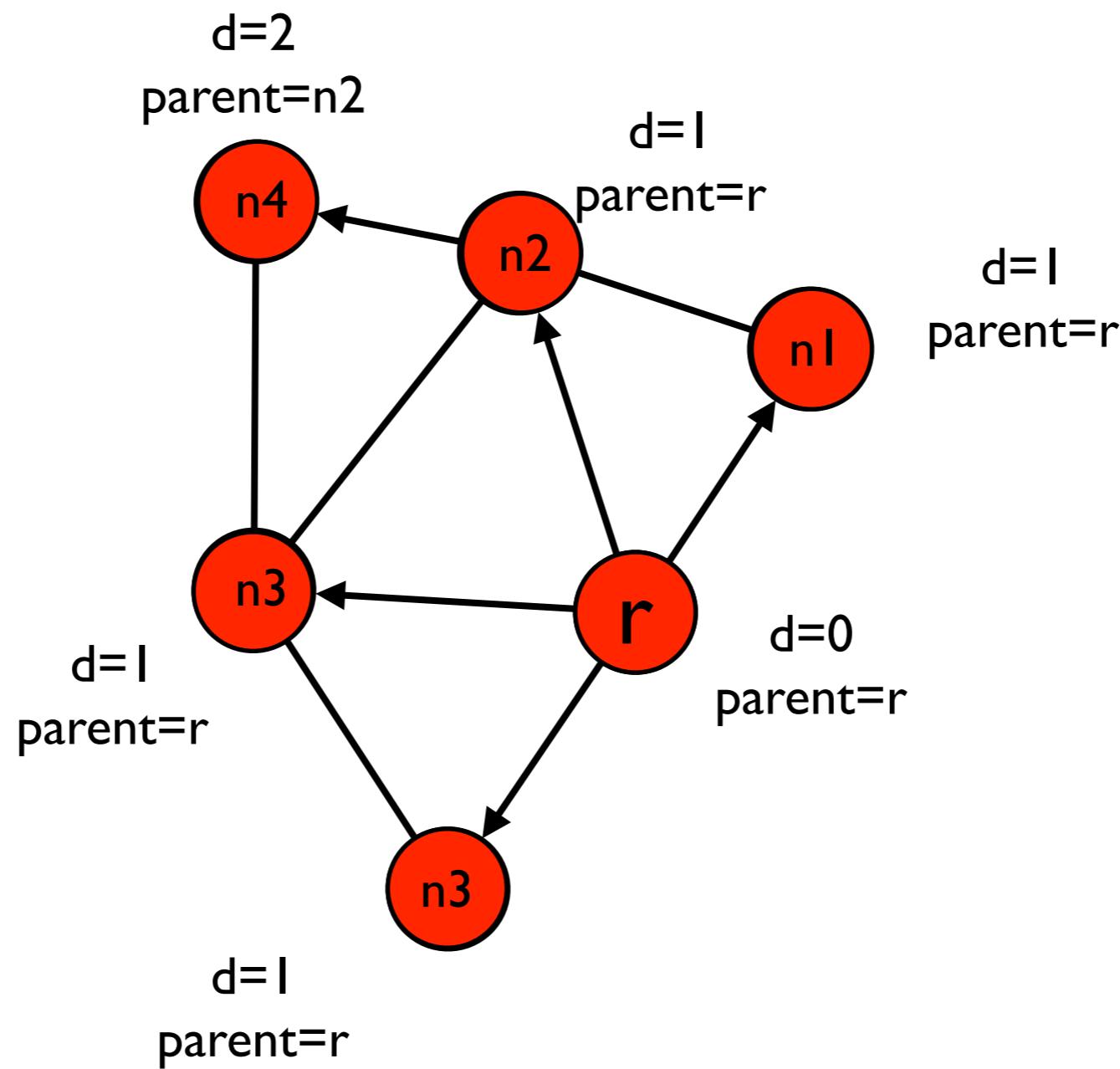
each node
broadcasts its d to
all its neighbors
except parent



Bellman-Ford-like BFS



Bellman-Ford-like BFS



Bellman-Ford-like BFS - Complexity

- Time Complexity: $O(\text{Diam}(G))$
- Message Complexity: $O(|E|)$

BFS - Complexity

	Time	Messages
Dijkstra-like	D^2	$E+nD$
Bellman-Ford-like	D	nE
Awerbuch-Peleg	$D\log^3 n$	$E+n\log^3 n$
Lower Bound	D	E

Spanning Tree and Routing

Routing

- «link state» (Dijkstra like)
 - Open Shortest Path First (OSPF)
- «distance vector» (Bellman-Ford like)
 - Border Gateway Protocol (BGP)
 - Interior Gateway Routing Protocol (IGRP, CISCO property)